

15-441/641: Computer Networks

Virtual Circuits, MPLS, VLAN

15-441 Fall 2019
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Fall 2019
<https://computer-networks.github.io/fa19/>

**Carnegie
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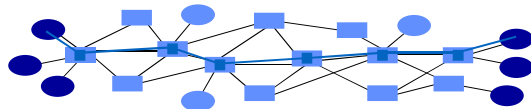
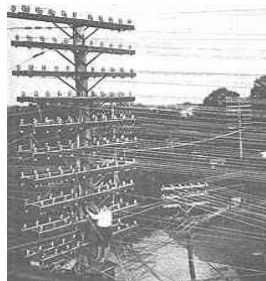
Outline

- Circuit switching refresher
- Virtual Circuits
 - Why virtual circuits?
 - How do they work?
- Today's virtual circuits: MPLS
- Virtual LANs
 - How do they differ?

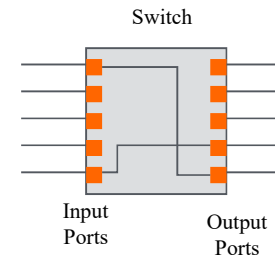


Circuit Switching

- Source first establishes a connection (circuit) to the destination.
 - Each router or switch along the way may reserve some bandwidth for the data flow
- Source sends the data over the circuit.
 - No destination address needed - routers know the path
- The connection is torn down.
- Example: traditional telephone network.



Circuit Switching



- Switches remembers how to forward data
 - No packets or addresses!
- Many options for switches
 - Connect specific wires (circuit = wire)
 - Forward on specific wire in specific timeslots (TDMA on each wire)
 - Forward to specific frequency on a specific wire (FDMA on each wire)



Circuit Versus Packet Switching

Circuit Switching

- Fast switches can be built relatively inexpensively
- Inefficient for bursty data
- Predictable performance (e.g. hard QoS)
- Requires circuit establishment before communication

Packet Switching

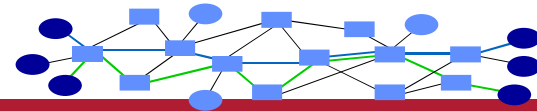
- Switch design is more complex and expensive
- Allows statistical multiplexing
- Difficult to provide QoS guarantees
- Data can be sent without signaling delay and overhead

Can we get the benefits of both?



Virtual Circuits

- Each wire carries many "virtual" circuits
- Forwarding based on virtual circuit (VC) identifier in a packet header
 - IP header: source IP, destination IP, etc.
 - Virtual circuit header: VC ID, ...
 - A path through the network is set up when the VC is established
 - Statistical multiplexing for efficiency, similar to IP
- Can support wide range of quality of service
 - No guarantees: best effort service
 - Weak guarantees: delay < 300 msec, ...
 - Strong guarantees: e.g. equivalent of physical circuit

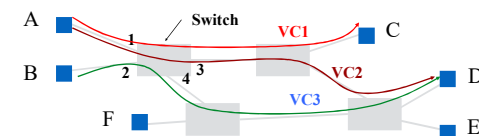


Virtual Circuits Versus Packet Switching

- Many similarities:
 - Forwarding based on "address" (VCID or destination address)
 - Statistical multiplexing for efficiency
 - Must have buffers space on switches
- Virtual circuit switching:
 - Uses short connection identifiers to forward packets
 - Switches maintain state for each connection so they can more easily implement features such as quality of service
 - Switches are stateful: VC connection state cannot be lost
- Packet switching:
 - Uses full destination addresses for forwarding packets
 - Can send data right away: no need to establish a connection first
 - Switches are stateless: easier to recover from failures
 - Adding QoS is hard, kind of – see QoS lecture



Virtual Circuit Forwarding

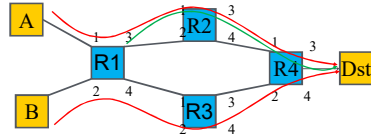


Address	Next Hop
VC1	3
VC2	3
VC3	4
VC4	?
VC5	?

- The address used for look up in the forwarding table is a virtual circuit identifier (VC id)
- Forwarding table entries are filled in during signaling
- VC id is often shorter than destination address and is typically "flat" (e.g., no CIDR)



VC versus Packets: Control over the End-to-End Path



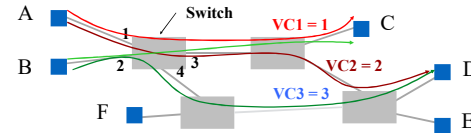
Different paths to same destination!
(useful for traffic engineering!)

R1 VC table:
VC 1 R2
VC 2 R3

R1 packet forwarding table:
Dst R2



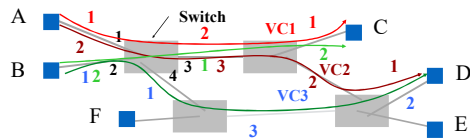
How to Pick a New VC Id?



- When B establishes green virtual circuit, how does it know what VC ids are available at all hops along the path?
- Even worse: every VC id may already be used on a link along the path to the destination!
- Solution: VC id swapping



VC id Swapping



Address	Next Hop	Next id
VC1=1	3	2
VC2=2	3	3
VC3=1	4	1
VC4=2	3	1

- Look up is based on VC id in header + incoming port number
- Forwarding table specifies outgoing port and new VC id
- VC id conflicts can be resolved locally during signaling

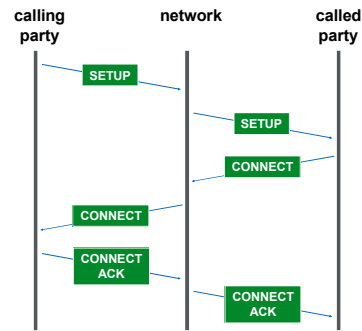


Connections and Signaling

- Permanent vs. switched virtual connections (PVC/SVC)
 - static vs. dynamic. PVCs last "a long time"
 - E.g., connect two bank locations with a PVC
 - SVCs are more like a phone call
- PVCs administratively configured (but not "manually")
- SVCs dynamically set up on a "per-call" basis
- Topology
 - point to point, point to multipoint, multipoint to multipoint
- Challenges: How to configure these things?
 - What VCI to use?
 - Setting up the path



Basic SVC Connection Setup



Virtual Circuits In Practice

- Asynchronous Transfer Mode - ATM: Teleco approach
 - Based on 53 byte "cells", not packets (I am not kidding)
 - Kitchen sink: design driven by voice requirements, but supports file transfer, video, etc.,
 - Intended as IP replacement. That didn't happen. :)
 - Today: dead.
- MPLS: The "IP Heads" answer to ATM
 - Stole all the good ideas from ATM and integrated them into IP
 - Today: Used inside many transit networks to provide traffic engineering, VPN support, ..
- Other networks just run IP.
- Older (ancient?) technology: Frame Relay
 - Only provided PVCs. Used for quasi-dedicated 56k/T1 links between offices, etc. Slower, less flexible than ATM.



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MPLS

- Multi-Protocol Label Switching
- Brings the virtual circuit concept into IP
- Driven by multiple forces
 - QoS, traffic engineering
 - Simplifies packet forwarding
- MPLS is implemented using an MPLS that sits between the IP and datalink header
 - VC ID is called an MPLS label

Layer 3 (IP) header

Layer 2 header

Layer 3 (IP) header

MPLS label

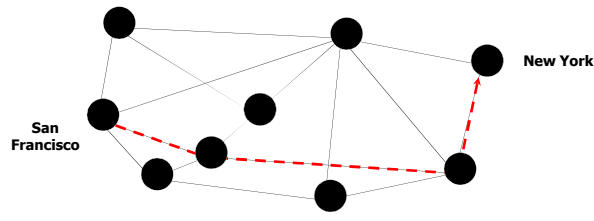
Layer 2 header

Some MPLS slides from H. Zhang



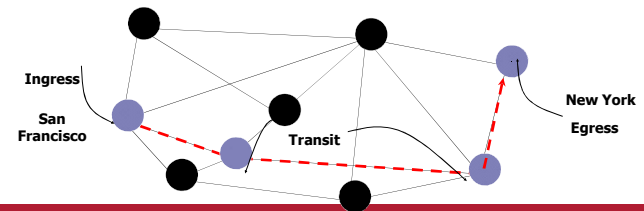
Label Switched Paths (LSP)

- Uni-directional path between two routers in an ISP's network
- Forces packet along a specific path (set of routers)



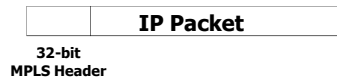
Label Switched Router (LSR)

- Performs LSP setup and MPLS packet forwarding
- Label Edge Router (LER): LSP ingress or egress
- Transit Router: forwards packet and swaps MPLS label



MPLS Header

- IP packet is encapsulated in MPLS header by the LSP ingress router and sent down a LSP



- MPLS forward packet based on the label
- IP packet is restored at end of LSP by egress router
 - TTL is adjusted, transit LSP routers count towards the TTL
 - MPLS is an optimization – it does not affect IP semantics



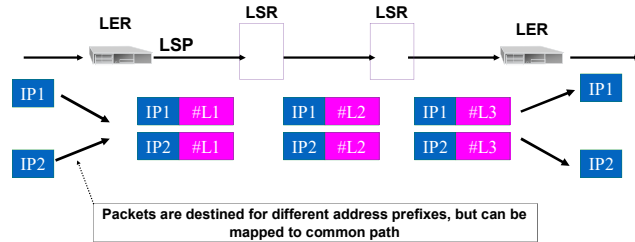
MPLS Header



- Label – 20 bits that identify LSP
- Class of service
- Stacking bit
 - Packets can be encapsulated in multiple MPLS headers
- Time to live
 - Decrement at each LSR, or
 - Pass through unchanged



Forwarding Equivalence Classes



- FEC = "A subset of packets that are all treated the same way by a LSR"
- The concept of FECs provides for a great deal of flexibility and scalability
- Can be used to force flows of different "sizes" (e.g., Mbps) to follow certain paths through the network – more flexible than traditional routing

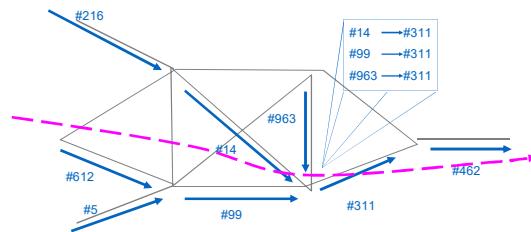


Establishing LSPs

- Use the Label Distribution Protocol (LDP) to establish paths based on IP forwarding tables
- Simple: IP routing protocols map forwarding of IP prefixes to LSPs as they fill in IP forwarding tables
- Establish new LSPs as needed
- MPLS packets follow the same path as IP – lose some MPLS benefits
- Explicitly establish LSPs to control flow of traffic
 - More work
 - Provides finer grain control over how traffic is distributed throughout the network
- Important tool for traffic engineering



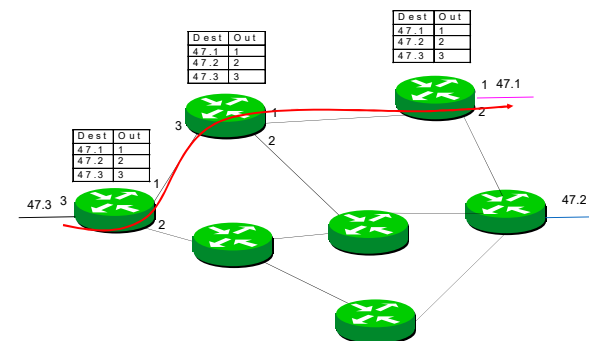
LSPs Driven by IP Routing



- A LSP is actually part of a sink tree with paths from every source to a destination (unidirectional).
- A control protocol (e.g. Label Distribution Protocol, LDP) builds the tree based on the IP forwarding tables.



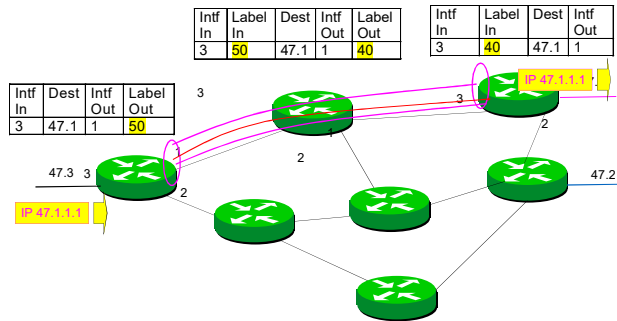
MPLS Builds on Standard IP



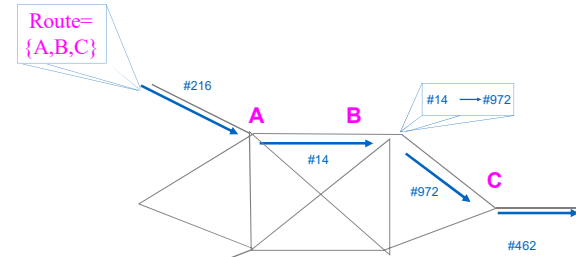
- Destination based forwarding tables as built by OSPF, IS-IS, RIP, etc.



Label Switched Path (LSP)

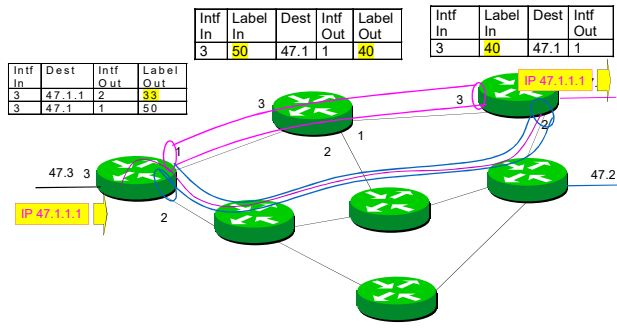


Explicitly Routed - ER-LSP



ER-LSP follows a route that is explicitly selected by a network manager.

Explicitly Routed LSP - Example



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VLAN Introduction

- VLANs logically segment switched LANs
 - Separates hardware topology from LAN topology
 - They operate at layer 2 (very different from MPLS!)
 - Partitioning is based on organization or function
 - It is independent of the physical location of nodes in the network
- Devices on a VLAN share their own (private) LAN
 - It is indistinguishable from a physical LAN, e.g., Ethernet, that has its own dedicated hardware (switches, wires)
 - Has all the same properties, e.g., broadcast capability
 - Form their own IP subnet



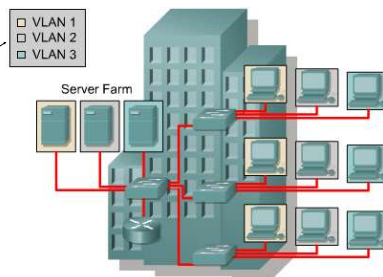
VLAN Benefits

- ◆ Performance: limits broadcast messages to the VLAN – improves scalability
 - ◆ E.g., very large organizations
 - ◆ Support for mobility in WiFi
- ◆ Management: manage network topology without changing the physical topology
 - ◆ E.g., departments in a university or company
- ◆ Security: isolates VLAN – VLANs connected by routers with smarter filtering capabilities
 - ◆ E.g., separate “guest” network from internal network so traffic is fully isolated



VLAN Example

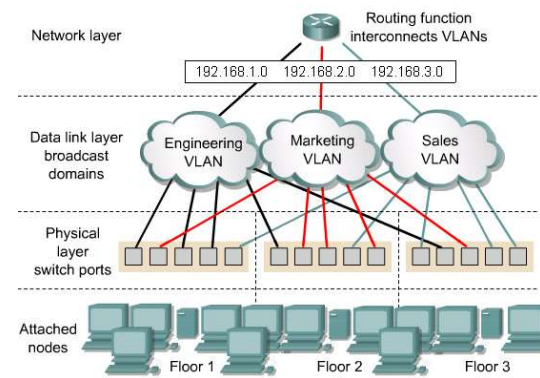
Devices with the same color form their own VLAN sharing the red physical hardware



- A switch creates a broadcast domain
- VLANs help manage broadcast domains
- VLANs can be defined on port groups, users, or protocols
- LAN switches and network management software provide a mechanism to create VLANs



VLAN Logical Topology



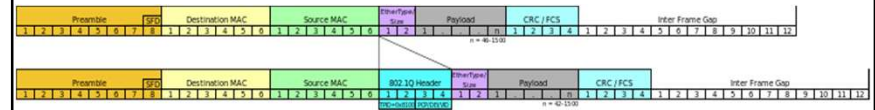
VLAN Types

- The VLAN for each packet is selected by a switch, not the host
- First switch adds a VLAN ID to the packet
 - Why?
- Last switch in the path removes the VLAN ID
- Add field to existing header or encapsulation
- VLAN memberships can be controlled in different ways, based on:
 - Port: incoming switch ports are tagged with VLAN ID
 - MAC address: switch has (MAC, VLAN ID) table
 - Protocol: switch as (protocol, VLAN ID) table



Example: 802.1Q Standard for VLANs over Ethernet

- A 32 bit VLAN header is inserted after the MAC addresses



- Header consists of
 - Tag Protocol Identifier (16b): single value that marks frame as a VLAN frame
 - Control bits (4b): mostly priority
 - VLAN Identifier (12b): identifies VLAN



Take Home Points

- Costs/benefits/goals of virtual circuits
- Tag/label swapping - basis for most VCs.
 - Makes label assignment link-local. Understand mechanism.
- MPLS - IP meets virtual circuits (links)
 - Used for VPNs, traffic engineering, reduced core routing table sizes
 - Management of ISPs at layer 3
- Virtual LANs – manage LANs in software
 - Simplifies management of edge networks at layer 2
 - Very widely used, e.g., cmu-guest versus cmu-secure WiFi access
 - Set up by manager based on organizational structure – no tag swapping

